

## 5.3 The MOV Attack

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# Introduction

- ✉ The MOV attack, named after Menezes, Okamoto, and Vanstone (1993), uses the Weil pairing to convert discrete log problem in  $E(\mathbb{F}_q)$  to one in  $\mathbb{F}_{q^m}^*$
- ✉ Let  $E$  be an elliptic curve over  $\mathbb{F}_q$ . Let  $P, Q \in E(\mathbb{F}_q)$ . Let  $N$  be the order of  $P$ . Assume that

$$\gcd(N, q) = 1.$$

Want to find  $k$  such that  $Q = kP$ .

# Lemma 5.1

## Lemma 5.1

$P, Q \in E(F_q)$ ,  $N = \text{ord}(P)$ ,  $\gcd(N, q) = 1$

$\exists k$  such that  $Q = kP$  if and only if  $NQ = \infty$  and  $e_N(P, Q) = 1$

# Proof of Lemma 5.1

Proof:

⊠ "only if" part:

$$Q = kP \rightarrow NQ = kNP = \infty$$

also  $e_N(P, Q) = e_N(P, P)^k = 1^k = 1$

⊠ "if" part:

$\gcd(N, q) = 1$ , we have  $E[N] \simeq \mathbb{Z}_N \oplus \mathbb{Z}_N$

Choose a point  $R$  such that  $\{P, R\}$  is a basis of  $E[N]$

Then  $Q = aP + bR$

By Corollary 3.10,  $e_N(P, R) = \zeta$  is a primitive  $N$ th root of unity.

If  $e_N(P, Q) = 1$ , then  $1 = e_N(P, Q) = e_N(P, P)^a e_N(P, R)^b = \zeta^b$

$$\rightarrow b \equiv 0 \pmod{N}$$

$$\rightarrow bR = \infty$$

$$\rightarrow Q = aP$$

# MOV attack - (1)

- ✉ Choose  $m$ , such that  $E[N] \subset E(F_{q^m})$   
By Corollary 3.11,  $\mu_N \subset F_{q^m}$
- ✉ MOV attack (Want to solve  $Q = kP$  for  $k$ .)
  - 1 Choose a random point  $T \in E(F_{q^m})$
  - 2 Compute  $M = \text{ord}(T)$
  - 3 Let  $d = \text{gcd}(M, N)$ , and let  $T_1 = (M/d)T$   
Then  $d = \text{ord}(T_1)$ ,  $d \mid N$ , so  $T_1 \in E[N]$
  - 4 Compute  $\zeta_1 = e_N(P, T_1)$  and  $\zeta_2 = e_N(Q, T_1)$   
Then both  $\zeta_1, \zeta_2 \in \mu_d \subseteq F_{q^m}^*$
  - 5 Solve discrete logarithm problem

$$\zeta_2 = \zeta_1^k \quad \text{in } F_{q^m}^*$$

This will give  $k \pmod{d}$

- 6 Repeat with random points  $T$  until the lcm of  $d$ 's is  $N$ .  
This determines  $k \pmod{N}$

## MOV attack - (2)

- ✉ Potentially,  $m$  could be large, in which case the discrete log problem in group  $F_{q^m}^*$  is just as hard as the original discrete log problem in  $E(F_q)$ . However, for supersingular curves, we can usually take  $m = 2$ .

## Proposition 5.3

### Proposition 5.3

$E/F_q$ , suppose  $a = q + 1 - \#E(F_q) = 0$

If  $\exists P \in E(F_q)$  of order  $N$ , then  $E[N] \subset E(F_{q^2})$

Proof:

$$(1) \quad \phi_q^2 - a\phi_q + q = 0$$

$$\rightarrow \phi_q^2 = -q \quad (\because a = 0)$$

$$(2) \quad \text{Let } S \in E[N]$$

$$\because \text{ord}(P) = N, \#E(F_q) = q + 1$$

$$\rightarrow N \mid q + 1 \quad \Leftrightarrow \quad -q \equiv 1 \pmod{N}$$

By (1),

$$\phi_q^2(S) = -qS = 1 \cdot S = S$$

By Lemma 4.5,  $S \in E(F_{q^2})$