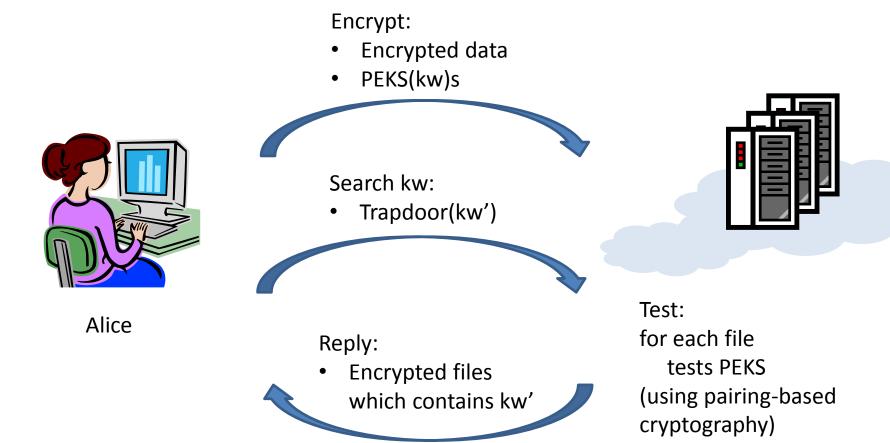
Searchable Encryption

陳榮傑

交通大學資工系 Cryptanalysis Lab 2012/03/29

Searchable Encryption

kw: keyword



Server gains no knowledge about kw or the file content stored on the Cloud Storage

Searchable Encryption

Encrypt:

• Encrypted data
• PEKS(kw)s

Search kw:
• Trapdoor(kw')

Test:
for each file

Server gains no knowledge about kw or the file content stored on the Cloud Storage

Encrypted files

which contains kw'

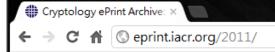
kw: keyword

tests PEKS

cryptography)

(using pairing-based

Data Example



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Johannes Blömer and Peter Günther and Gennadij Liske

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Differential Attacks on Generalized Feistel Schemes

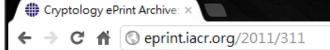
Valerie Nachef and Emmanuel Volte and Jacques Patarin

2011/704 (PS PS.GZ PDF)

Security Analysis of a PUF based RFID Authentication Protocol

Masoumeh Safkhani and Nasour Bagheri and Majid Naderi

Data Example



Cryptology ePrint Archive: Report 201 (/311)

Targeted Malleability: Homomorphic Encryption for Restricted Computations

Dan Boneh and Gil Segev and Brent Waters

Abstract: We put forward the notion of targeted malleability: given a homomorphic encryption scheme, in variou computations one can perform on encrypted data. We introduce a precise framework, generalizing the foundation Naor (SICOMP '00), ensuring that the malleability of a scheme is targeted only at a specific set of "allowable" fur

In this setting we are mainly interested in the efficiency of such schemes as a function of the number of repeated he ciphertext grows linearly with the number of such operations is straightforward, obtaining more realistic (or merely

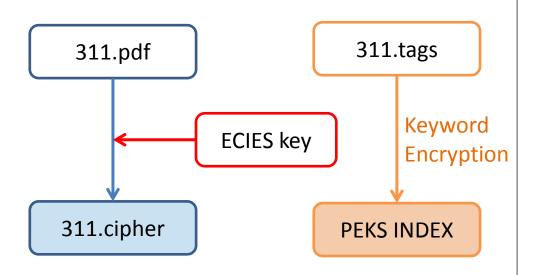
We present two constructions that transform any homomorphic encryption scheme into one that offers targeted mand on succinct non-interactive arguments, which are currently known to exist in the standard model based on vacconstructions offer somewhat different efficiency guarantees, each of which may be preferable depending on the united transformations.

Categor Keywords: foundations / Homomorphic encryption, non-malleable encryption

id: 311
title: Targeted Malleability: Homomorphic Encryption for Restricted Computations author: Dan Boneh author: Gil Segev author: Brent Waters keyword: foundations keyword: Homomorphic encryption keyword: non-malleable encryption

Keyword Lists (311.tags)

Public Key Encryption with Keyword Search (PEKS) (1/3)



311.pdf

Targeted Malleability: Homomorphic Encryption for Restricted Computations

Dan Boneh*

Gil Segev[†]

Brent Waters[‡]

Abstract

We put forward the notion of targeted malleability: given a homomorphic encryption scheme, in various scenarios we would like to restrict the homomorphic computations one can perform on encrypted data. We introduce a precise framework, generalizing the foundational notion of non-malleability introduced by Dolev, Dwork, and Naor (SICOMP '00), ensuring that the malleability of a scheme is targeted only at a specific set of "allowable" functions.

In this setting we are mainly interested in the efficiency of such schemes as a function of the number of repeated homomorphic operations. Whereas constructing a scheme whose ciphertext grows linearly with the number of such operations is straightforward, obtaining more realistic (or merely non-trivial) length guarantees is significantly more challenging.

We present two constructions that transform any homomorphic encryption scheme into one that offers targeted malleability. Our constructions rely on standard cryptographic tools and on succinct non-interactive arguments, which are currently known to exist in the standard model based on variants of the knowledge-of-exponent assumption. The two constructions offer somewhat different efficiency guarantees, each of which may be preferable depending on the underlying building blocks.

Keywords: Homomorphic encryption, Non-malleable encryption.

^{*}Stanford University. Supported by NSF, DARPA, and AFOSR.

[†]Microsoft Research, Mountain View, CA 94043, USA.

University of Texas at Austin. Supported by NSF CNS-0716199, CNS-0915361, and CNS-0952692. DARPA PROCEED, Air Force Office of Scientific Research (AFO SR) MURI, DHS Grant 2006-CS-001-000001-02, and the Sloan Foundation.

Search Keyword

Search kw:

 Trapdoor(author: Brent Waters, Alice's private key)



Alice

Reply:

 Encrypted files which contains kw

311.aes128

311.rsa



Test:

for each file tests PEKS in

PEKS INDEX

Search Keyword

File 1: PEKS(id) Trapdoor(author: Brent Waters) PEKS(title) PEKS(author1) PEKS(author2) File 2: PEKS(id) Trapdoor(author: Brent Waters) PEKS(title) PEKS(author) PEKS(keyword1) PEKS(keyword2)

 $PEKS = (g^r, H_2(t)), \qquad t$

PEKS INDEX

$$t = e(H_1(KW), h^r),$$

$$h = g^{\alpha}$$

 $Trapdoor = H_1(KW)^{\alpha}$

Server tests each PEKS whether $H_2(e(H_1(KW)^{\alpha}, g^r) = H_2(t)$

Implementation

- PBC Library by Ben Lynn
- Tate Pairing
- supersingular curve: $y^2 = x^3 + x$ over F_q
 - embedding degree k = 2
 - q is a prime and $q \equiv 3 \mod 4$
 - q is 1536-bit long
 - group order r is 256-bit long

Key Length:

	Wor Neconinendations (2011)							. •
Date	Minimum of Strength	Symmetric Algorithms	Asymmetric	Discrete Key	Logarithm Group	Elliptique Curve	Hash (A)	Hash (B)
> 2030	128	AES-128	3072	256	3072	256	SHA-256 SHA-384 SHA-512	SHA-1 SHA-224 SHA-256 SHA-384 SHA-512

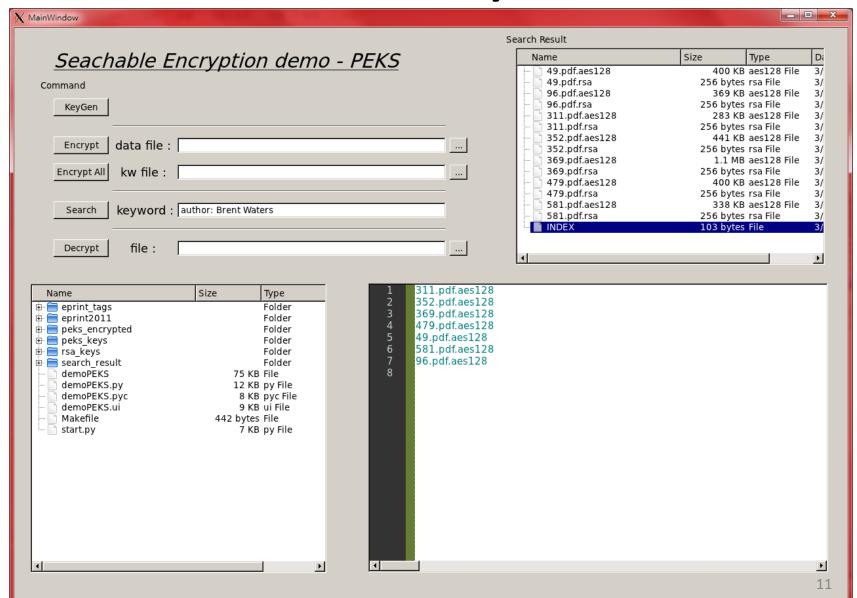
Recommendations (20

Forthcoming Research

To enhance search capability

```
Query Type
Equality query: (x_i = a)
                                       for any a \in T
Comparison query: (x_i \ge a) for any a \in T
Subset query: (x_i \in A)
                                      for any A \subseteq T
Equality conjunction: (x_1 = a_1) \wedge \ldots \wedge (x_w = a_w)
Comparison conjunction: (x_1 \geq a_1) \wedge \ldots \wedge (x_w \geq a_w)
Subset conjunction: (x_1 \in A_1) \land \ldots \land (x_w \in A_w)
```

Demo PEKS library



• Questions?

Thank you